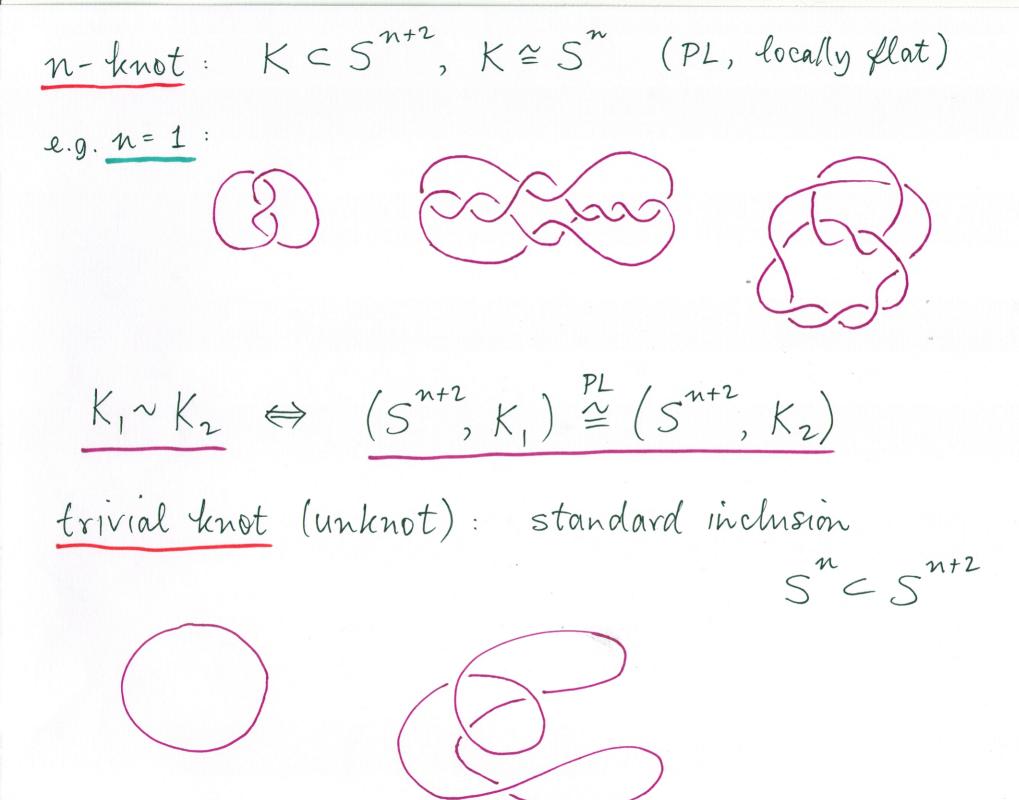
UNSOLVABLE PROBLEMS ABOUT HIGHER-DIMENSIONAL KNOT GROUPS Joint with

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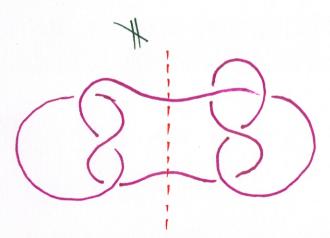
$$G(K) = \Pi_{1}(S^{n+2}-K)$$
, the group of K

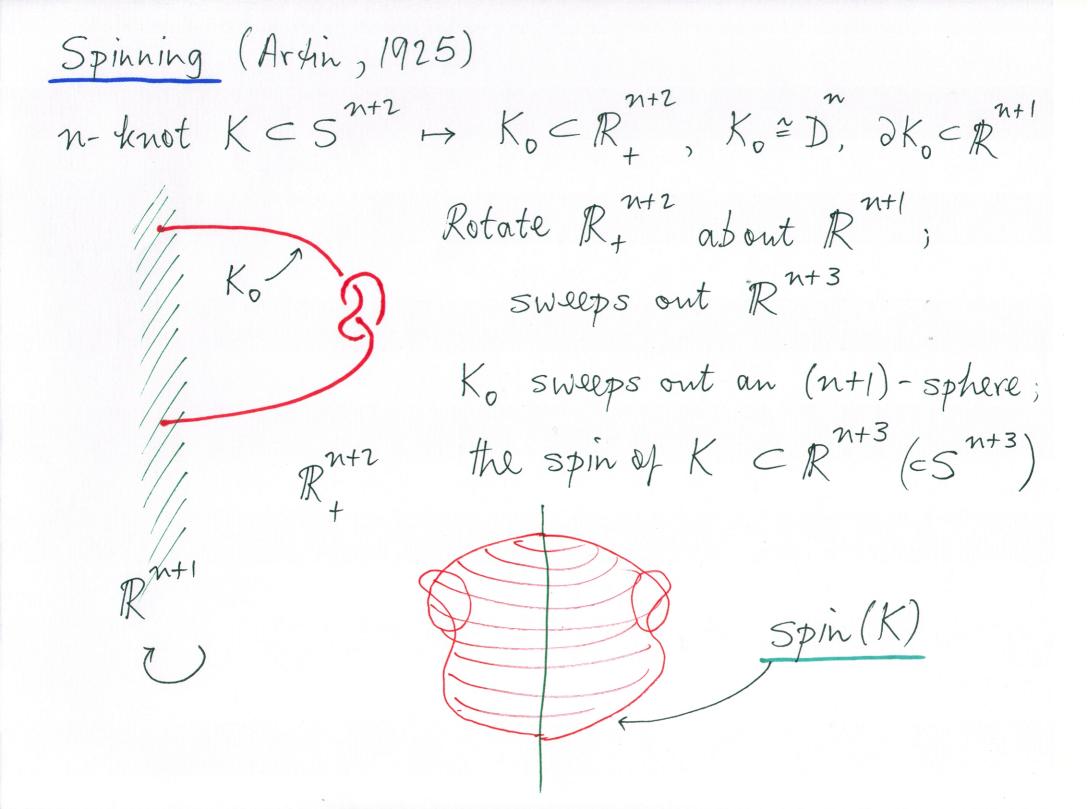
$$K_{1} \sim K_{2} \implies G(K_{1}) \cong G(K_{2})$$

$$G(unknot) \cong \mathbb{Z}$$

$$n=1:$$
 $G(K) \cong \mathbb{Z} \Rightarrow K = unknot$

$$K_1$$
 prime, $G(K_1) \cong G(K_2) \Rightarrow K_1 \sim K_2$





 $G(spin K) \cong G(K)$

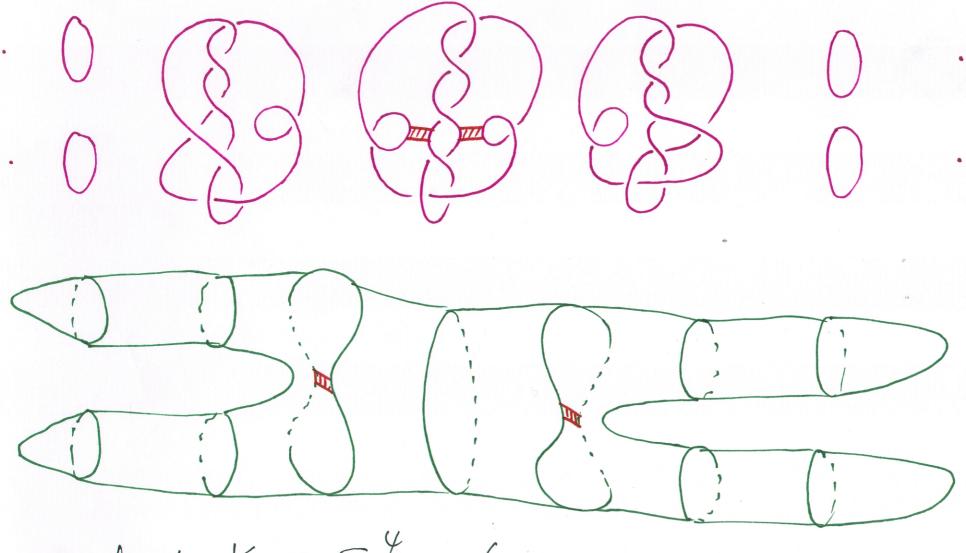
So 3 non-trivial n-knots 7n >1.

 $n=2: G(K) \cong \mathbb{Z} \Rightarrow K \text{ ToPologically} \sim \text{unknot}$ (Freedman, 1982)

Q. Is K PL~ unknot?

 \exists (TOP) prime 2-knots K_1, K_2 with $G(K_1) \cong G(K_2)$, $K_1 \not\sim K_2$

n73: I non-trivial n-knots with G(K) ≈ Z



$$G(K) = \langle a, t : t^{-1}at = a^{-1}, a^{3} = 1 \rangle$$

Let
$$K_{0n} = \{n-knot \text{ groups}\}$$
 $K_{00} = \{Z\}; \text{ spinning shows } K_{0n} \subset K_{n+1} \quad \forall n \geqslant 0.$

$$[G \text{ group}; \quad \forall \text{ complex with } \pi_{1}(Y) \cong G, \quad \pi_{n}(Y) = 0, \quad n \neq 1.$$

Then $H_{k}(G) = H_{k}(Y)$
 $K \subset S^{n+2}$ an $n-knot$
 $X = S^{n+2} - K$
 $G = G(K) = \pi_{1}(X)$

$$X = finite complex \Rightarrow G \in \mathcal{G} = \{finitely \text{ presented groups}\}$$
 $H_{*}(X) \cong H_{*}(S^{1})$
 $H_{1}(X) \cong \mathbb{Z} \Rightarrow H_{1}(G) = G/[G,G] \cong \mathbb{Z}$
 $H_{2}(X) = 0 \Rightarrow H_{3}(G) = 0$

$$H_{2}(X) = 0 \Rightarrow H_{2}(G) = 0$$

Van Kampen >

$$G/\langle S^{n+2}\rangle = \pi_{i}(S^{n+2}) = 1.$$

M= meridian

Thm (Kervaire, 1965) For
$$n73$$
, $G \in \mathcal{K}_{on}$ iff

(1) $G \in \mathcal{G}$; (2) $H_1(G) \cong \mathbb{Z}$; (3) $H_2(G) = 0$;

(4) G has weight 1.

50 $\{7\} = \mathcal{K}_0 \subset \mathcal{K}_1 \subset \mathcal{K}_2 \subset \mathcal{K}_3 = \mathcal{K}_{on}$, $n73$
 $\mathcal{K}_0 \neq \mathcal{K}_1$: (8) $G = \langle x, y : x^2 = y^3 \rangle \neq \mathbb{Z}$
 $\mathcal{K}_{o1} \neq \mathcal{K}_{o2}$: $\langle t, a : t^- at = a^2, a^3 = 1 \rangle \in \mathcal{K}_{o2} - \mathcal{K}_{o1}$

 $\mathcal{K}_{2} \neq \mathcal{K}_{3}$: $\langle t, \alpha : t | \alpha t = \alpha^{2}, \alpha^{5} = 1 \rangle \in \mathcal{K}_{2} - \mathcal{K}_{2}$

Q. Is there a "nice" characterization of Lez? $S_n = \{\pi, (S^{n+2} - N) : N \text{ conn., closed, orientable } n - mfld \}$ $Mn = \{ \pm 1/(M^{n+2} - N) : M \text{ closed}, 1 - conn. (n+2) - mpld; \}$ $S_0 = M_0 = \mathcal{X}_0$; $S_1 = M_1 = \mathcal{X}_1$ $S_2 = S_n, n_2 = S$ $M_2 = M_n$, N_72 (=M)So $S = \{ \pi, (S - N) : N \text{ conn. closed, orientable } \}$ surface }

Thu (Simon, 1980).
$$G \in S$$
 iff (1) $G \in G$; (2) $H_1(G) \cong Z$;

(3) $\exists t \in G$ such that $\langle t \rangle = G \times t \wedge C_t = H_2(G)$.

($C_t = \text{centrahizer of } t$; $\Lambda = \text{Pontyagain product}$)

 $G \in M$ iff (1) $G \in G$, \times (2) G has weight 1.

{ $\mathcal{I}_3 = \mathcal{K}_0 \subset \mathcal{K}_1 \subset \mathcal{K}_2 \subset \mathcal{K}_3 \subset \mathcal{S} \subset M \subset G$

All inclusions proper: $A_5 \times \mathcal{I} \in \mathcal{S} - \mathcal{K}_3$;

 $\langle x, y, s : x^2 = y^3 = (xy)^7, s^2 \times s = x^2 \rangle \in M - \mathcal{S}$
 $\mathcal{I}_{XZ} \in G - M$

Let B = Q = b, where Efinite presentations of groups e Q } is recursively unmerable. The recognition problem Rec (a, B) is solvable if there is an algorithm to deade, given a finite présentation of a group $A \in \mathbb{Q}$, whether or not $A \in \mathcal{B}$. Thm (Adjan, Radin, 1958)

Rec(b, {13) is unsolvable.

ie there is no algorithm to decide whether or not the group defined by a given finite presentation is trivial. Based on
Thu (Novikov, Boone, 1955-56). If Ely with
unsolvable word problem.

ie. There is no algorithm to decide, given a word w in a set of generators of G, whether or not $[w] = J \in G$.



Pyotr Sergeyevich Novikov 1901–1975



William Werner Boone 1920–1983



It has recently been announced from Russia that the 'word problem in groups' is not solvable. This is a decision problem not unlike the 'word problem in semi-groups', but very much more important, having applications in topology: attempts were being made to solve this decision problem before any such problems had been proved unsolvable. ...Another problem which mathematicians are very anxious to settle is known as 'the decision problem of the equivalence of manifolds'.... It is probably unsolvable, but has never been proved to be so. A similar decision problem which might well be unsolvable is the one concerning knots which has already been mentioned. (Turing, 1954)

Thm (Haken, 1902) I algorithm to decide whether or not a given 1-knot KC53 is trivial. Since K trivial $\Leftrightarrow G(K) \cong \mathbb{Z}$.

Cor. Rec (Ke1, Ko) is solvable.

G = { Ko, K1, K2, K3, S, M, y3

If $a \in G$ then Efinike presentations of groups $\in a$ is recursively enumerable.

 $\mathcal{T}_{hm}(Gonzalez-Acnña-G.-Simon).$ Let $Q,B \in G$, $\mathcal{B} \neq Q$, $Q > \mathcal{K}_{e_3}$. Then Rec(Q,B) is unsolvable.

Q. Does this also hold for Q = Kez?

Most cases proved by using unsolvability of Rec(G, E13). e.g.

Propⁿ. I effective procedure which takes a finite presentation of a group G and produces a finite presentation of a group H such that

(1) $H \in \mathcal{U}_3$; (2) $G = 1 \Rightarrow H \cong \mathbb{Z}$; (3) $G \neq 1 \Rightarrow H \notin \mathcal{U}_2$.

Cor. of LocB < Loz then Rec (K3, B) is unsolvable. $(a sol'n. to Rec(K_3, B) wd. \mapsto sol'n. to Rec(b, {1}).)$ Addendum. In Proph. Gembleds in H. Cor. I H & Kez with unsolvable word problem. Cor. IHE Les which contains a copy of every GEG. Q. Does every Gey embed in some HEKez?

n-knot triviality problem {n-knots} is recursively enumerable. So can ask: " given an n-knot, is it trivial?" Solvable of n=1. Thm (G-A-G-S.) If F ∈ Len with unsolvable word problem then the n-knot triviality problem is unsolvable. Cor (Nasutovsky-Weinslerger, 1996) If n73 the n-tenot triviality problem is unsolvable. Q. Is there a G E Kez with unsolvable word problem?

(5ⁿ⁺², K) n-hnot, n72, with Sketch proof. G = G(K) having unsolvable word problem. Surgery on $K: M^{n+2} = (5^{n+2} - N(K)) U(D^{n+1} \times S^1)$ $S^n \times D^2$ $C = *xS^1 = D^{n+1}xS^1 = M$ $\pi_1(M) \cong G$; [C] = m = meridian of KLet & be a trivial n-sphere < M $\pi_1(M-\Sigma_1) \cong G \times \mathbb{Z}(\sigma), \quad \sigma = \text{meridian } \emptyset \leq 1$

w word in the generators of G, [w] \in [6]. Tw loop in M-E such that [Tw] = o'w'own $\in \pi_{i}(M-S_{i})$ $[Y_w] = M \in \pi_1(M)$.. Tw 1sotopic to C in M Do surgery on The $S \mapsto n - sphel S_{1,r} \subset S^{n+2}$

$$[W] = 1 \in G : [Sw] = M \in \Pi_{1}(M-\Sigma)$$

$$\therefore Sw \text{ isotopic to } C \text{ in } M-\Sigma$$

$$\therefore (S^{n+2}, \Sigma_{w}) \text{ trivial}$$

$$[W] \neq 1 \in G : \Pi_{1}(S^{n+2} - \Sigma_{w}) = \langle G, \sigma : \sigma'w\sigma = w_{n} \rangle$$

$$= HNN \text{ extension of } G ; \therefore \neq \mathbb{Z}$$

$$\therefore (S^{n+2}, \Sigma_{w}) \text{ non-trivial}.$$

in a solution to the n-tenot triviality problem would give a solution to the word problem for G. It

simple connectivity of 3-manifolds problem: "given a connected, closed 3-nyld. M, is $\pi_1(M) = 1?$ " (*) Let $M_3 = \{ \pi_1(M) : M \text{ closed } 3 - \text{mylld.} \}$ { finite presentations of groups in M3 } is rec. enumerable. (an show if $G \in M_3$ then $G = 1 \Leftrightarrow H_1(G) = H_3(G) = 0$ (*) is solvable: (1) $\pi_1(M) = 1 \Leftrightarrow M \cong S^3$ (Perellman, 2003)

(2) I algorithm to decide whether or not a given 3-mfld. M is ≅ S³ (Rubinstein-Thompson, 1994)

Gey => H, (G), Hr (G) finishly generated abelian groups (Hy (G), 473, not nec. fin. gen. (Stallings, 1963)). clearly I algorithm to compute HIG) for Geg. Thm (6, 1980) There is no algorithm to decide, given $G \in \mathcal{G}$, whether or not $H_2(G) = 0$. More generally:

Thu (G-A-G-S). Let $I \subset \mathbb{N}$, $I \neq \emptyset$, $\{1\}$. Then "given $G \in \mathcal{G}$, is $H_k(G) = 0 \quad \forall k \in I$ " is unsolvable.