

# A topological dynamical system with two different positive sofic entropies

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# Overview

- Sofic groups
- Sofic entropy
- Main theorem
- Proper 2-colorings of random hyper-graphs

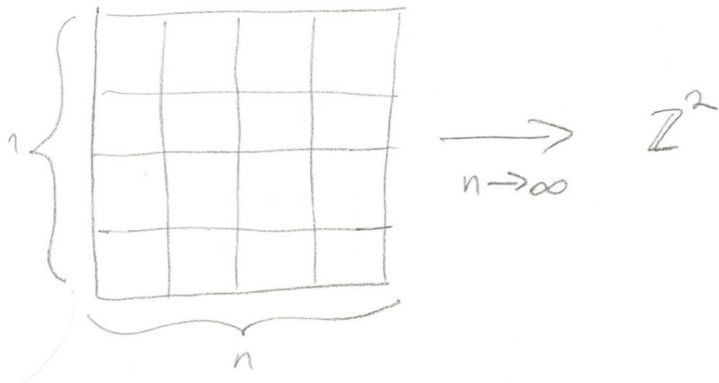
# Benjamini-Schramm convergence

## Definition

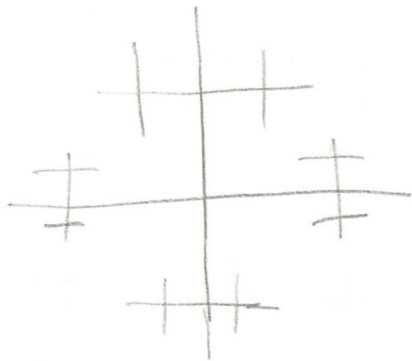
A sequence  $G_i = (V_i, E_i)$  of finite graphs **BS-converges** to a fixed vertex-transitive graph  $G = (V, E)$  if for every radius  $r > 0$ ,

$$\lim_{i \rightarrow \infty} \frac{\#\{v \in V_i : B_r(G_i, v) \cong B_r(G, o)\}}{\#V_i} = 1.$$

# The square grid as a BS-limit



?

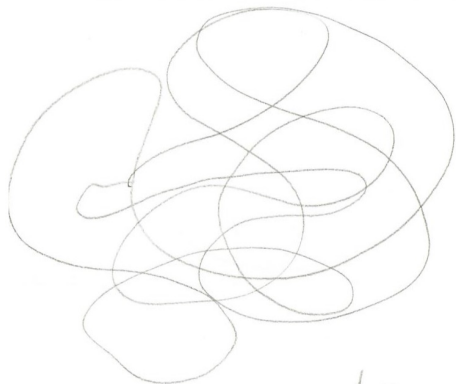


?

6

4-regular  
depth  $n$  tree

# 4-regular tree



random 4-regular  
graph on  
 $n$  vertices

$\xrightarrow{n \rightarrow \infty}$



4-reg Tree!

# Sofic groups

Let  $\Gamma$  be a countable group with finite generating set  $S$ .

The **Cayley graph** of  $(\Gamma, S)$  has vertex set  $V = \Gamma$  and edges

$$E = \{(g, gs) : g \in \Gamma, s \in S\}.$$

A **sofic approximation** to  $\Gamma$  is a sequence  $\Sigma = \{G_i = (V_i, E_i)\}$  of finite graphs that BS-converges to the Cayley graph.

$\Gamma$  is **sofic** if it has a sofic approximation.

## Question

Are all countable groups sofic?

(and now for something completely different)

Theorem (B.-Burton, 2019)

*For  $d \geq 5$ , either  $PSL(d, \mathbb{Z})$  is not flexibly stable or there exists a non-sofic group.*



## Sofic entropy (overview)

Given a continuous action  $\Gamma \curvearrowright X$ , the  $\Sigma$ -entropy (or **sofic entropy**) is

$$h_{\Sigma}(\Gamma \curvearrowright X) \in \{-\infty\} \cup [0, \infty].$$

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## Main properties

- 1 Monotone under embeddings:

$$Y \subset X \Rightarrow h_{\Sigma}(\Gamma \curvearrowright Y) \leq h_{\Sigma}(\Gamma \curvearrowright X).$$

- 2  $h_{\Sigma}(\Gamma \curvearrowright A^{\Gamma}) = \log |A|$ .
- 3  $\Gamma$  amenable  $\Rightarrow h_{\Sigma}(\Gamma \curvearrowright X) =$  classical topological entropy (Kerr-Li, 2012).
- 4 There is a measure entropy version (B. 2010) and a variational principle (Kerr-Li, 2011).

# Symbolic dynamics

Let

- $A$  be a finite ‘alphabet’,
- $A^\Gamma = \{x : \Gamma \rightarrow A\}$ ,
- $\Gamma \curvearrowright A^\Gamma$  by  $(g \cdot x)_f = x_{g^{-1}f}$ .
- $X \subset A^\Gamma$  be closed and  $\Gamma$ -invariant.  $X$  is a **subshift**.

## The entropy of a subshift

Let  $\Sigma = \{G_n = (V_n, E_n)\}$  be a sofic approximation.

A coloring  $\chi : V_n \rightarrow A$  is an  $(r, \epsilon)$ -**model** for  $\Gamma \curvearrowright X$  on  $G_n$  if there exists a subset  $W_n \subset V_n$  such that

- $|W_n| \geq (1 - \epsilon)|V_n|$ ,
- for all  $v \in W_n$ , there exists  $x \in X$  such that  $B_r(G_n, \chi, v) \cong B_r(\Gamma, x, 1_\Gamma)$ .

### Definition

$$h_\Sigma(\Gamma \curvearrowright X) := \inf_{r>0} \inf_{\epsilon>0} \limsup_{n \rightarrow \infty} |V_n|^{-1} \log \#\{(r, \epsilon)\text{-models for } \Gamma \curvearrowright X \text{ on } G_n\}.$$

### Question

How does  $\Sigma$ -entropy depend on  $\Sigma$ ?

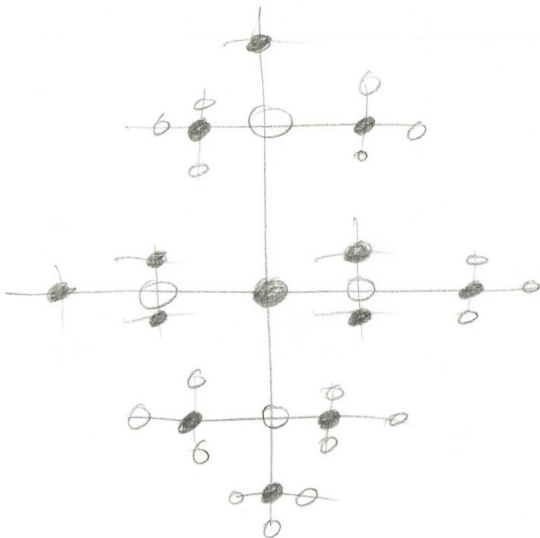
# A degenerate example

Let

- $\Gamma = \langle a, b \rangle$ ,
- $X = \{x, y\} \subset \{0, 1\}^\Gamma$  be the two bi-partite colorings of  $\Gamma$ ,
- $\Sigma = \{G_n\}$  be a sofic approximation to the Cayley graph.

- If each  $G_n$  is bi-partite then  $h_\Sigma(\Gamma \curvearrowright X) = 0$ .
- If  $G_n$  is 'random' then  $h_\Sigma(\Gamma \curvearrowright X) = -\infty$ .

# A degenerate example



# The main result

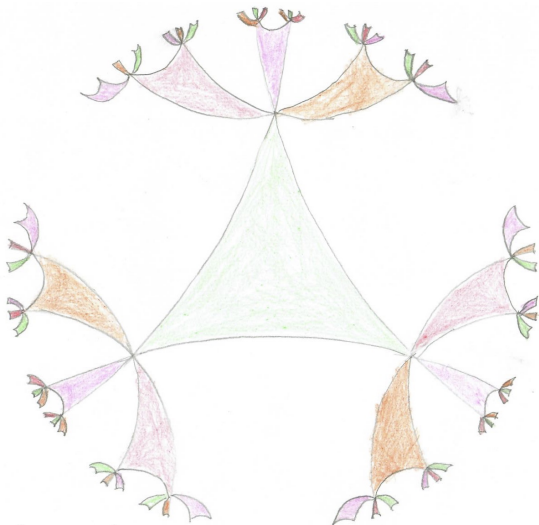
## Theorem (Airey-B.-Lin)

*There exists an explicit group action  $\Gamma \curvearrowright X$  and sofic approximations  $\Sigma_1, \Sigma_2$  such that*

$$0 < h_{\Sigma_1}(\Gamma \curvearrowright X) < h_{\Sigma_2}(\Gamma \curvearrowright X).$$

The measure entropy case is still open.

# The example



$$\Gamma = \langle s_1, \dots, s_d : s_1^k = s_2^k = \dots = s_d^k = 1 \rangle.$$



# Hyper-graphs and proper colorings

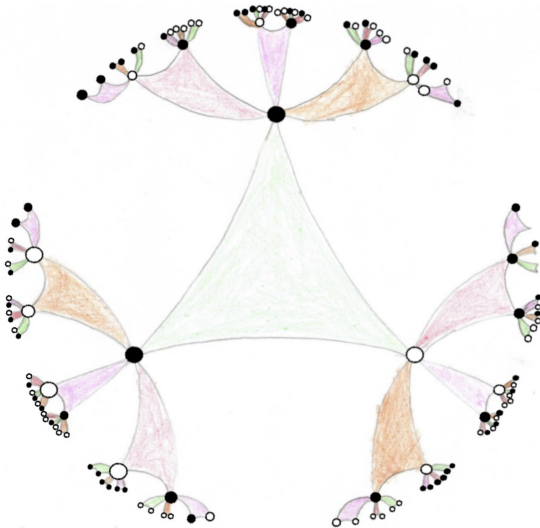
A **hyper-graph** is a pair  $G = (V, E)$  where  $E$  is a collection of subsets of  $V$ .

A coloring  $\chi : V \rightarrow \{0, 1\}$  is **proper** if  $\chi(e) = \{0, 1\}$  for every  $e \in E$ .

The **Cayley hyper-graph** of  $\Gamma$  has  $V = \Gamma$  and  $E = \{\text{left-cosets of generator subgroups}\}$ .

Let  $\mathbf{X} \subset \{0, 1\}^\Gamma$  be the subset of proper colorings of the Cayley hyper-graph.

# A proper coloring



# Sofic approximations

Fix a sofic approximation  $\Sigma = \{G_n = (V_n, E_n)\}$  to  $\Gamma$ .

A coloring  $\chi : V_n \rightarrow \{0, 1\}$  is  **$\epsilon$ -proper** if

$$\#\{\mathbf{e} \in E_n : |\chi(\mathbf{e})| = 1\} \leq \epsilon n.$$

If  $\Sigma$  is any sofic approximation then

$$h_\Sigma(\Gamma \curvearrowright X) = \inf_{\epsilon > 0} \limsup_{n \rightarrow \infty} n^{-1} \log \#\{\epsilon\text{-proper colorings of } G_n\}.$$

# Random sofic approximations

Suppose  $V_n$  is a sequence of finite sets and  $\mathbb{P}_n$  are probability measures on the set of hyper-graphs on  $V_n$ . Then  $\mathbb{P}_n$  is a **random sofic approximation** if for every  $r > 0$  and  $\delta > 0$  there exists  $\epsilon > 0$  such that

$$\mathbb{P}_n \left( \frac{\#\{v \in V_n : B_r(G_n, v) \cong B_r(G, o)\}}{\#V_n} > 1 - \delta \right) > 1 - n^{-\epsilon n}$$

for all large enough  $n$ .

## Main observation

If  $\mathbb{P}_n$  is a random sofic approximation and  $\Omega_n$  are collections of hyper-graphs with  $\mathbb{P}_n(\Omega_n) \geq e^{-cn}$  for some  $c > 0$  then  $\exists$  a sofic approximation  $\Sigma = \{G_n\}$  with  $G_n \in \Omega_n$  for all  $n$ .

# The uniform model

Given a homomorphism  $\sigma : \Gamma \rightarrow \text{sym}(n)$ , let  $G_\sigma$  be the hyper-graph on  $[n]$  with edges being the orbits of the generators:

$$E_\sigma = \left\{ \left\{ \sigma(s_i^j)v \right\}_{j=0}^{k-1} : i \in [d], v \in [n] \right\}.$$

$\sigma$  is a **uniform** if  $G_\sigma$  is  **$k$ -uniform**.

The **uniform model** is the uniform probability measure  $\mathbb{P}_n^{\text{unif}}$  on the set of uniform homomorphisms. By abuse, we also say the **uniform model** is the induced sequence of distributions on hyper-graphs. It is a random sofic approximation.

# Equitable colorings

Let

- $Z(\epsilon; \sigma) := \#\{\epsilon\text{-proper colorings of } G_\sigma\}$ .
- $Z(\sigma) := \#\{\text{proper colorings of } G_\sigma\}$ ,
- a coloring  $\chi : V \rightarrow \{0, 1\}$  is **equitable** if  $|\chi^{-1}(0)| = |\chi^{-1}(1)| = |V|/2$ .
- $Z_e(\sigma) := \#\{\text{equitable proper colorings of } G_\sigma\}$ .

Then

$$\begin{aligned}\mathbb{E}_n^{\text{unif}}[Z(\epsilon; \sigma)] &\approx e^{O(\epsilon)n} \mathbb{E}_n^{\text{unif}}[Z(\sigma)] \\ \mathbb{E}_n^{\text{unif}}[Z(\sigma)] &\approx \mathbb{E}_n^{\text{unif}}[Z_e(\sigma)].\end{aligned}$$

## The satisfiability threshold

Fix a large  $k$  and let  $r = d/k$ . Conjecturally, there exists a satisfiability threshold  $r_{\text{sat}} > 0$  such that

- if  $r < r_{\text{sat}}$  then w.h.p, proper colorings exist,
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 Dimitris Achlioptas and Cristopher Moore.

Random  $k$ -SAT: two moments suffice to cross a sharp threshold,  
*SIAM J. Comput.*, 36(3):740–762, 2006.

 Amin Coja-Oghlan and Lenka Zdeborová.

The condensation transition in random hypergraph 2-coloring,  
*arXiv:1107.2341*, 35 pages, 2011.

 Amin Coja-Oghlan and Lenka Zdeborová.

The condensation transition in random hypergraph 2-coloring,  
*Proceedings of the Twenty-Third Annual ACM-SIAM Symposium on Discrete Algorithms*, 241–250. ACM, New York, 2012.



# A tale of two moments

Let

- $r_{\text{first}}$  := the value of  $r$  with  $\lim_{n \rightarrow \infty} n^{-1} \log \mathbb{E}_n^{\text{unif}}[Z_e(\sigma)] = 0$ .
- $r_{\text{second}}$  := the largest value of  $r$  with

$$\lim_{n \rightarrow \infty} n^{-1} \log \mathbb{E}_n^{\text{unif}}[Z_e^2(\sigma)] = \lim_{n \rightarrow \infty} n^{-1} \log \mathbb{E}_n^{\text{unif}}[Z_e(\sigma)]^2.$$

Then

$$r_{\text{sat}} \in [r_{\text{second}}, r_{\text{first}}].$$

Theorem (cf. Achlioptas-Moore, 2006)

$$\begin{aligned} r_{\text{first}} &= \frac{\log(2)}{2} 2^k - \frac{\log(2)}{2} + O(2^{-k}) \\ r_{\text{second}} &= \frac{\log(2)}{2} 2^k - \frac{\log(2) + 1}{2} + O(2^{-k}). \end{aligned}$$

# The second moment

## Lemma

$\mathbb{E}_n^{\text{unif}}[Z_e^2] = \mathbb{E}_n^{\text{unif}}[Z_e] \mathbb{E}_n^{\text{unif}}[Z_e | \chi \text{ is proper}]$  where  $\chi$  is any equitable coloring.

Let  $\mathcal{H} = \text{Hom}_{k\text{-uniform}}(\Gamma, \text{sym}(n))$ .

$$\begin{aligned} \mathbb{E}_n^{\text{unif}}[Z_e^2] &= |\mathcal{H}|^{-1} \sum_{\sigma \in \mathcal{H}} \left( \sum_{\chi \text{ equitable}} \mathbf{1}_{(\chi \text{ is proper w.r.t. } G_\sigma)} \right)^2 \\ &= |\mathcal{H}|^{-1} \sum_{\sigma \in \mathcal{H}} \sum_{\chi, \tilde{\chi} \text{ equitable}} \mathbf{1}_{(\chi \text{ is proper w.r.t. } G_\sigma)} \mathbf{1}_{(\tilde{\chi} \text{ is proper w.r.t. } G_\sigma)} \\ &= \sum_{\chi, \tilde{\chi} \text{ equitable}} \mathbb{P}_n^{\text{unif}}(\chi \text{ is proper and } \tilde{\chi} \text{ is proper}) \\ &= \sum_{\chi, \tilde{\chi} \text{ equitable}} \mathbb{P}_n^{\text{unif}}(\chi \text{ is proper}) \mathbb{P}_n^{\text{unif}}(\tilde{\chi} \text{ is proper} | \chi \text{ is proper}) \\ &= \mathbb{E}_n^{\text{unif}}[Z_e] \mathbb{E}_n^{\text{unif}}[Z_e | \chi \text{ is proper}]. \end{aligned}$$

# The planted model

The **planted model** is  $\mathbb{P}_n^{\text{plant}}(\cdot) = \mathbb{P}_n^{\text{unif}}(\cdot | \chi \text{ is proper})$  where  $\chi$  is any equitable coloring. It is a random sofic approximation.

- $r < r_{\text{second}} \Rightarrow \mathbb{E}_n^{\text{unif}}[Z_e] \approx \mathbb{E}_n^{\text{plant}}[Z_e],$
- $r_{\text{second}} < r \Rightarrow \mathbb{E}_n^{\text{unif}}[Z_e] \ll \mathbb{E}_n^{\text{plant}}[Z_e].$

# The *enhanced* second moment

Theorem (cf. Coja-Oghlan-Zdeberova, 2011)

For  $r \in \left( r_{\text{second}}, r_{\text{second}} + \frac{1 - \log(2)}{2} \right)$ ,

$$\mathbb{P}_n^{\text{plant}} \left( Z_e \leq \mathbb{E}_n^{\text{unif}}[Z_e] \right) = \Omega(1/n^C).$$

Therefore,

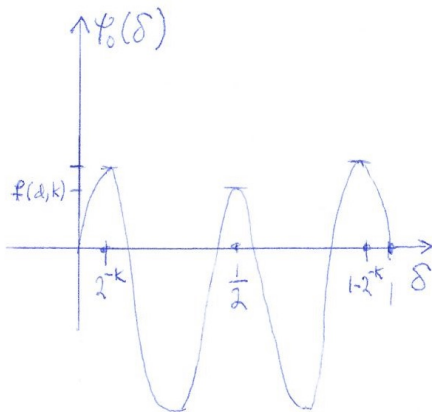
- $\mathbb{P}_n^{\text{unif}} \left( Z_e \approx \mathbb{E}_n^{\text{unif}}[Z_e] \right) = \Omega(1/n^C)$ ,
- $\mathbb{P}_n^{\text{plant}} \left( Z_e \gg \mathbb{E}_n^{\text{unif}}[Z_e] \right) = \Omega(e^{-cn})$ ,
- $r_{\text{second}} < r_{\text{sat}}$ ,
- $\exists \Sigma_1, \Sigma_2$  such that

$$0 < h_{\Sigma_1}(\Gamma \curvearrowright X) = f(d, k) < h_{\Sigma_2}(\Gamma \curvearrowright X).$$

# Expected number of proper colorings at distance $\delta$ from the planted colorings

## Theorem

Let  $Z_e(\delta)$  be the number of proper colorings at Hamming distance  $\delta n$  from the planted coloring. Then  $\mathbb{E}_n^{plant}[Z_d(\delta)] \approx e^{n\psi_0(\delta)}$  where ...



# The local cluster

Fix an equitable  $\chi : [n] \rightarrow \{0, 1\}$ . This is the **planted coloring**.

The **local cluster** is the set of  $G_\sigma$ -proper equitable colorings  $\psi : [n] \rightarrow \{0, 1\}$  such that  $d_{\text{Hamming}}(\chi, \psi) \leq 2^{-k/2}n$ .

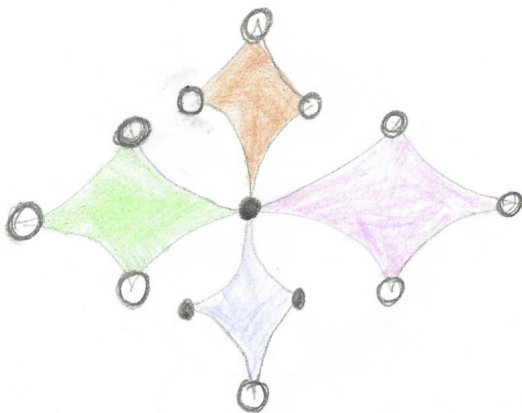
In order to prove

$$\mathbb{P}_n^{\text{plant}} (Z_e \leq \mathbb{E}_n^{\text{unif}}[Z_e]) = \Omega(1/n^C),$$

it suffices to show

$$\mathbb{P}_n^{\text{plant}} (|\text{local cluster}| \leq \mathbb{E}_n^{\text{unif}}[Z_e]) = \Omega(1/n^C).$$

## A supporting role



A vertex  $v$  **supports** an edge  $e$  if  $\chi(v) \cap \chi(e \setminus \{v\}) = \emptyset$ .

# The core

Let  $C_0 = [n]$ , and  $C_{l+1}$  be the set of all  $v \in [n]$  such that there are at least three edges  $e$  such that

- $v$  supports  $e$ ,
- $e \setminus \{v\} \subset C_l$ .

Note  $C_0 \supset C_1 \supset \dots$ .

The **core** is  $C_\infty = \bigcap_l C_l$ .



# The rigid core

## Lemma

With high probability (and a few lies),

- $\frac{|C_\infty|}{n} = 1 - 2^{-k}(1 + o_k(1))$ .
- If  $\psi : [n] \rightarrow \{0, 1\}$  is a proper coloring then either
  - ▶  $\psi \upharpoonright C_\infty = \chi \upharpoonright C_\infty$  or
  - ▶  $\#\{v \in [n] : \psi(v) \neq \chi(v)\} \geq \text{constant} \cdot \frac{n}{k^4}$ .

$\Rightarrow$  whp  $\#\{\psi : d(\psi, \chi) \leq 2^{-k/2}\} \leq 2^{n-|C_\infty|} \leq \mathbb{E}_n^{\text{unif}}[Z_e]$ .

# Open problems

- Provide general necessary or sufficient conditions (on subshifts of finite type say) for when the sofic entropy range is greater than a singleton.
- The measure entropy case is **wide** open.
- Is there a well-defined notion of mod  $p$   $\ell^2$ -Betti numbers that does not depend on the choice of sofic approximation?

Thank you!