

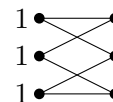
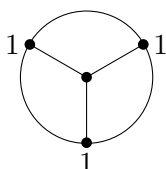
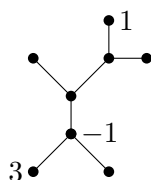
# PROBLEM SET 1 - DHAR'S BURNING ALGORITHM AND RIEMANN-ROCH

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The graphs in this problem set are all finite, but notice that metric graphs would only require minor changes.

### Warm up questions.

- (1) Compute the rank of the following divisors.



- (2) Let  $G$  be the banana graph consisting of two vertices and  $g + 1$  edges between them. Let  $D$  be the divisor with  $a$  chips on one of the vertices and  $b$  chips on the other, where  $0 \leq a \leq b$ . Show that

$$r(D) = \begin{cases} a & b \leq g \\ a + b - g & \text{otherwise} \end{cases}$$



- (3) Let  $D$  be a  $v$ -reduced divisor such that  $D(v) = a$ . Show that  $r(D) \leq a$ . Find an example where  $r(D)$  is strictly smaller than  $a$ .
- (4) Let  $G$  be a trivalent graph of genus  $g$  with  $E$  edges and  $V$  vertices.
- Express  $E$  in terms of  $g$ .
  - Let  $D$  be a divisor of degree  $E \cdot V$ . Show that there must be a vertex  $v$ , whose valency is smaller or equal  $D(v)$ . Conclude that every reduced divisor of degree  $E \cdot V$  is effective.
  - Prove that every divisor of degree  $E \cdot V + r$  has rank at least  $r$ .
- (5) Use Riemann–Roch to show that the rank of every divisor of degree  $d > 2g - 2$  is  $d - g$ . Show that a divisor of degree  $2g - 2$  has rank  $g - 1$  if and only if it is equivalent to the canonical divisor.

**Dhar's burning algorithm.** Let  $D$  be a divisor on a finite graph  $G$ , and fix a vertex  $q$  on the graph.

- (1) Show that  $D$  is equivalent to a divisor that is effective everywhere away from  $q$  (but may be very negative at  $q$ ).

If the divisor is not reduced, then by definition there is a set of vertices which does not contain  $q$ , such that firing from them does not create a pole. Fire them, and check whether we get a reduced divisor.

- (2) Show that after repeating this finitely many times, we end up with a reduced divisor.
- (3) Let  $A$  be the set where a rational function  $f$  obtains its minimum, and let  $\text{div}(A)$  the divisor obtained by firing each of the vertices of  $A$  once. Show that  $\text{div}(f)(v) \leq \text{div}(A)(v)$  at every vertex  $v$  of  $A$ . Conclude that every divisor is equivalent to a *unique* reduced divisor.

**Riemann-Roch.** The following notations and definitions will be useful throughout:

- For a rational function  $f$ , we have  $\text{div}(f)(v) = \sum (f(v) - f(u))$ , where the sum is taken over the vertices adjacent to  $v$ . In other words, the divisor  $\text{div}(f)$  is obtained by chip-firing  $-f(v)$  times from each vertex  $v$ .
- For a set  $A$  of vertices, the divisor  $\text{div}(A)$  is the principal divisor obtained by firing once from each vertex of  $A$ .
- Given an orientation  $\mathcal{O}$ , there is an associated divisor  $D_{\mathcal{O}}$  whose degree at every vertex is the number of incoming edges minus 1. A divisor of the form  $D_{\mathcal{O}}$  is called a *moderator*.
- A *source* is a vertex for whom all the adjacent edges are oriented outwards.
- An orientation is *acyclic* if it doesn't have any oriented cycles.
- Given a divisor  $D$ , we define  $\text{deg}^+(D)$  as the sum of the non-negative coefficients of  $D$ .

- (1) Show that for an orientation  $\mathcal{O}$ , the divisor  $D_{\mathcal{O}}$  has degree  $g - 1$ .
- (2) Show that  $D_{\mathcal{O}} + D_{\mathcal{O}^-}$  is equal to the canonical divisor, where  $\mathcal{O}^-$  represents the reverse orientation of  $\mathcal{O}$ .

From now on, we fix an acyclic orientation  $\mathcal{O}$ .

- (3) Show that  $D_{\mathcal{O}}$  is not equivalent to an effective divisor:
  - Assume by contradiction that  $D_{\mathcal{O}} + \text{div}(f)$  is effective, and let  $A$  be a connected component of the set on which  $f$  obtains its minimum.
  - Show that  $(D_{\mathcal{O}} + \text{div}(f))|_A \leq (D_{\mathcal{O}} + \text{div}(A))|_A$ .
  - Show that the restriction of  $\mathcal{O}$  to  $A$  has a source.
  - Conclude that  $D_{\mathcal{O}} + \text{div}(f)$  is not effective.

Let  $D$  be a  $q$ -reduced divisor with respect to some vertex  $q$ . When starting a fire from  $q$  (in the sense of the burning algorithm) it burns through the whole graph. At each vertex, choose a direction in which the fire spreads. Fixing such a choice determines an acyclic orientation  $\mathcal{O}$ , and a moderator  $\nu = D_{\mathcal{O}}$ .

- (4) Show that  $D(v) \leq D_{\mathcal{O}}(v)$  for every  $v \neq q$ , and conclude that either  $D$  or  $\nu - D$  is effective.

- (5) Show that for every moderator  $\nu$  associated to an orientation obtained as above, the rank of  $D$  is strictly smaller than  $\deg^+(D - \nu)$ .
- (6) In fact, show that the rank of  $D$  equals  $m - 1$ , where  $m$  is the minimum of  $\min(\deg^+ D' - \nu)$  as  $\nu$  varies over all the moderators, and  $D'$  varies over the divisor class of  $D$ .
- (7) Prove the Riemann–Roch theorem, namely,

$$r(D) - r(K - D) = \deg(D) - g + 1.$$